DSAA 5012 Advanced Data Management for Data Science

LECTURE 9 EXERCISES RELATIONAL DATABASE DESIGN: FUNCTIONAL DEPENDENCIES





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Assume that this table contains the *only* set of tuples that may appear in a relation R(X, Y, V, W). Which of the following FDs hold in R?

tuple	Х	Y	V	W
1	x ₁	У ₁	V ₁	W ₁
2	x ₁	У ₁	V ₂	W ₂
3	x ₂	У ₁	V ₁	W ₃
4	x ₂	У ₁	V ₃	W ₄

X→X	Yes	- trivial (holds in any relation) $4 + x_2 + y_1 + y_1 + y_2 + y_2 + y_2 + y_1 + y_2 + y_2 + y_2 + y_1 + y_2 + y_$
Х→Ү	Yes	 – for a given X value all Y values are identical
X→V	No	- V values differ for same X value (e.g., tuples 1 & 2)
X→W	No	- W values differ for same X value (e.g., tuples 1 & 2)
Y→X	No	- X values differ for same Y value (e.g., tuples 2 & 3)
W→X	Yes	 – all W values are different
XV→Y	Yes	– X alone determines Y
YV→X	No	– X values differ for same YV value in tuples 1 and 3



In Exercise 1, we assumed that we know all possible records in the table, which is not usually true. In general, by looking at an instance of a relation, we can only tell FDs that are <u>not</u> satisfied.

List 5 FDs that are <u>not</u> satisfied in the table.

А	В	С
a ₁	b ₁	С ₁
a ₁	b ₁	С ₂
a ₂	b ₁	С ₁
a ₂	b ₁	С ₃

A—	γС
В—	→A
В—	ъС
C—	→A
AB-	→C



Given relation schema R(X, Y, U, V, W) and $F = \{X \rightarrow Y, UV \rightarrow W, V \rightarrow X\}$

a) Determine the closure of each attribute.

 $X^{+} = \{X, Y\}$ (Look for X on LHS of FDs) $Y^{+} = \{Y\}$ $U^{+} = \{U\}$ $V^{+} = \{V, X, Y\}$ $W^{+} = \{W\}$

b) What are the candidate keys of R?

The candidate key is UV since $UV^{+} = \{X, Y, U, V, W\}$.





Given relation schema R(A, B, C, G, H, I) and $F = \{A \rightarrow B, \qquad A \rightarrow C, \qquad CG \rightarrow H, \qquad CG \rightarrow H, \qquad B \rightarrow H\}$ a) Is AG a (super)key of R given F? Yes Why? Since $AG \rightarrow R$ Compute AG⁺ $AG^{(0)} = \{A, G\}$ $AG^{(1)} = \{A, G, B\}$ $(A \rightarrow B \text{ and } A \subset \{A, G\})$ $\begin{array}{ll} \mathsf{A}\mathsf{G}^{(2)} = \{\mathsf{A},\,\mathsf{G},\,\mathsf{B},\,\mathsf{C}\} & (\mathsf{A}{\rightarrow}\mathsf{C} \text{ and }\mathsf{A}{\subseteq}\{\mathsf{A},\,\mathsf{G}\}) \\ \mathsf{A}\mathsf{G}^{(3)} = \{\mathsf{A},\,\mathsf{G},\,\mathsf{B},\,\mathsf{C},\,\mathsf{H}\} & (\mathsf{C}\mathsf{G}{\rightarrow}\mathsf{H} \text{ and }\mathsf{C}\mathsf{G}{\subseteq}\{\mathsf{A},\,\mathsf{G},\,\mathsf{B},\,\mathsf{C}\}) \end{array}$ (CG \rightarrow I and CG \subseteq {A, G, B, C, H}) $AG^{(4)} = \{A, G, B, C, H, I\}$ b) Is AG a candidate key? Yes Why? Cannot remove A or G c) Does $A^+ \rightarrow R$ hold? No since $A^+ = \{A, B, C, H\}$ d) Does $G^+ \rightarrow R$ hold? No since $G^+ = \{G\}$





Complete Exercises 5, 6 and 7; upload your completed worksheet to Canvas by 11 p.m.





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Given relation schema R(A, B, C, D, E) and $F = \{A \rightarrow B, AB \rightarrow C, D \rightarrow AC\}$

a) Determine the following attribute closures.

 $A^{+}= \{A, B, C\} \qquad B^{+}= \{B\} \qquad C^{+}= \{C\} \qquad D^{+}= \{D, A, B, C\} \qquad E^{+}= \{E\}$

c) Find a canonical cover of *F*.

 $F_c = \{A \rightarrow B, AB \rightarrow C, D \rightarrow AC\}$

 $F_{c} = \{A \rightarrow B, A \rightarrow C, D \rightarrow AC\}$ $F_{c} = \{A \rightarrow BC, D \rightarrow AC\}$ $F_{c} = \{A \rightarrow BC, D \rightarrow A\}$

B extraneous in AB \rightarrow C since A \rightarrow AB & AB \rightarrow C \Rightarrow A \rightarrow C use union rule C extraneous in D \rightarrow AC since D \rightarrow A & A \rightarrow C \Rightarrow D \rightarrow C



We want to create the database for a bank that contains accounts (A), branches (B) and customers (C).

- a) What are the functional dependencies implied by the following constraints?
 - i. An account cannot be shared by multiple customers. Account \rightarrow Customer $A \rightarrow C$
 - ii. Two different branches do not have the same account. Account \rightarrow Branch $A \rightarrow B$
 - iii. Each customer can have at most one account in a branch (but different accounts in different branches).

Branch, Customer \rightarrow Account BC \rightarrow A

b) What are the candidate keys?

(Branch, Customer) and Account

BC and A



Given relation schema R(A, B, C). Assume we do not know the keys of the relation. Write a <u>valid</u> SQL query whose result can be used to determine if A is a potential candidate key. Explain how to interpret the query result to determine if A is a potential candidate key.



Write an SQL query to test if the FD $A \rightarrow B$ holds in R?

 Same as the first query above, but replace the last line with having count(distinct B)>1

